

Shortest Paths

- Shortest Paths
- Properties of Shortest Paths
- Dijkstra's Algorithm
- Shortest Paths on DAGs

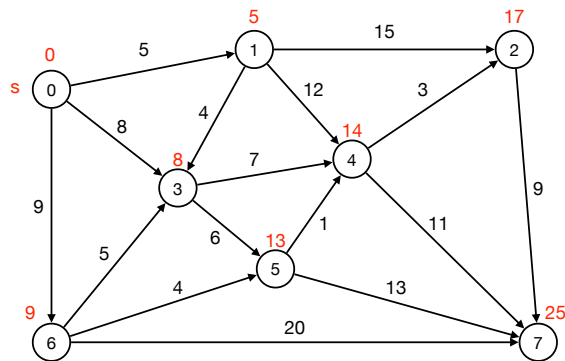
Philip Bille

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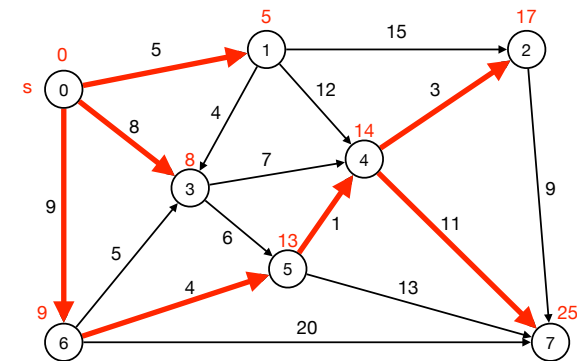
Shortest Paths

- **Shortest paths.** Given a directed, weighted graph G and vertex s , find shortest path from s to all vertices in G .



Shortest Paths

- **Shortest paths.** Given a directed, weighted graph G and vertex s , find shortest path from s to all vertices in G .
- **Shortest path tree.** Represent shortest paths in a tree from s .



Applications

- Routing, scheduling, pipelining, ...

Shortest Paths

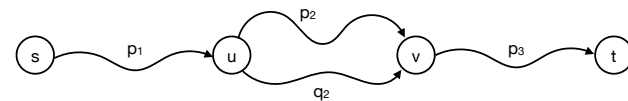
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- **Properties of Shortest Paths**
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Properties of Shortest Paths

- **Assume for simplicity:**
 - All vertices are reachable from s .
- \implies a (shortest) path to each vertex always exists.

Properties of Shortest Paths

- **Subpath property.** Any subpath of a shortest path is a shortest path.
- **Proof.**
 - Consider shortest path from s to t consisting of p_1 , p_2 and p_3 .



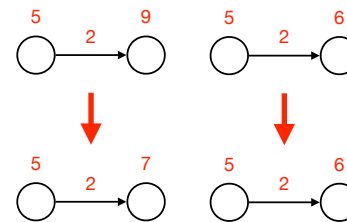
- Assume q_2 is shorter than p_2 .
- \implies Then p_1 , q_2 and p_3 is shorter than p .

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Dijkstra's Algorithm

- **Goal.** Given a directed, weighted graph with **non-negative weights** and a vertex s , compute shortest paths from s to all vertices.
- **Dijkstra's algorithm.**
 - Maintains **distance estimate** $v.d$ for each vertex v = length of shortest **known** path from s to v .
 - Updates distance estimates by **relaxing** edges.

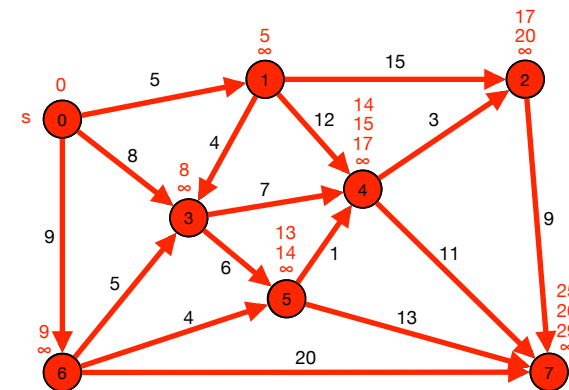
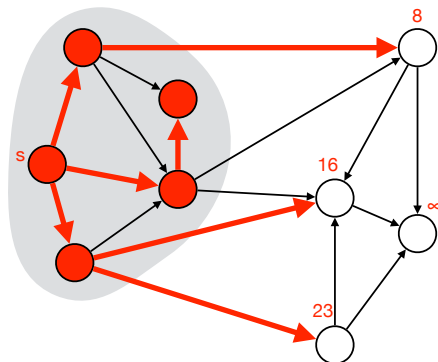


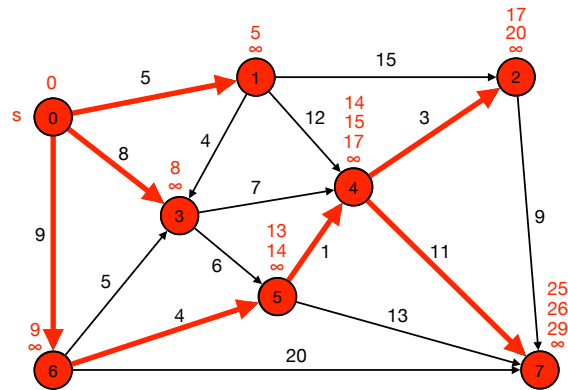
```

RELAX(u,v)
  if (v.d > u.d + w(u,v))
    v.d = u.d + w(u,v)
  
```

Dijkstra's Algorithm

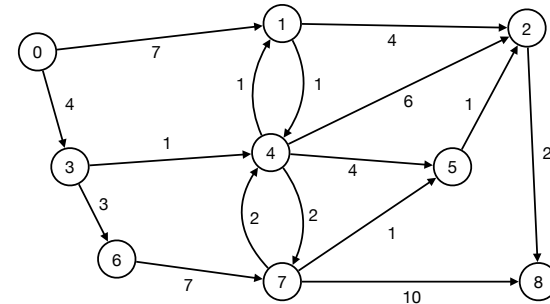
- Initialize $s.d = 0$ and $v.d = \infty$ for all vertices $v \in V \setminus \{s\}$.
- Grow tree T from s .
- In each step, add vertex with **smallest** distance estimate to T .
- Relax all outgoing edges of v .





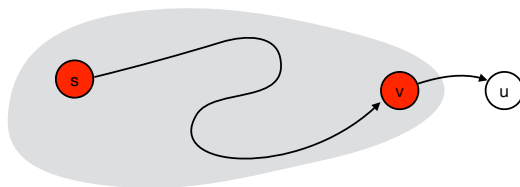
Dijkstra's Algorithm

- Initialize $s.d = 0$ and $v.d = \infty$ for all vertices $v \in V \setminus \{s\}$.
- Grow tree T from s .
- In each step, add vertex with **smallest** distance estimate to T .
- Relax all outgoing edges of v .
- **Exercise.** Show execution of Dijkstra's algorithm from vertex 0.



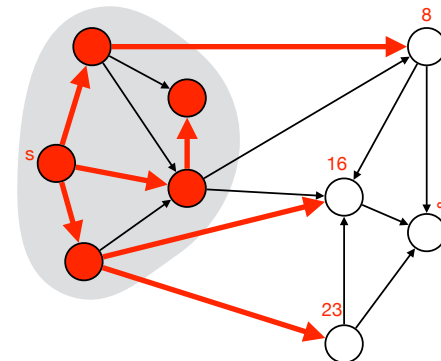
Dijkstra's Algorithm

- **Lemma.** Dijkstra's algorithms computes shortest paths.
- **Proof.**
 - Consider some step after growing tree T and assume distances in T are correct.
 - Consider closest vertex u of s **not** in T .
 - Shortest path from s to u ends with an edge $e = (v,u)$.
 - v is closer than u to $s \implies v$ is in T . (u was **closest** not in T)
 - \implies shortest path to u is in T except e .
 - e is relaxed \implies distance estimate to v is correct shortest distance.
 - Dijkstra adds e to $T \implies T$ is shortest path tree after $n-1$ steps.



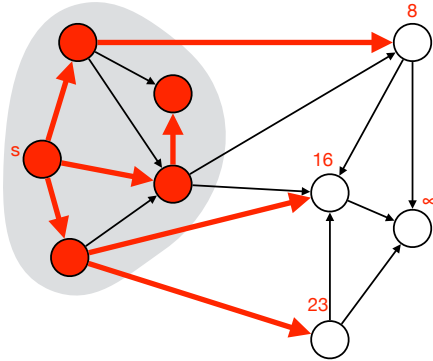
Dijkstra's Algorithm

- **Implementation.** How do we implement Dijkstra's algorithm?
- **Challenge.** Find vertex with smallest distance estimate.



Dijkstra's Algorithm

- **Implementation.** Maintain vertices outside T in priority queue.
 - **Key** of vertex $v = v.d$.
 - In each step:
 - Find vertex u with smallest distance estimate = EXTRACT-MIN
 - Relax edges that u point to with DECREASE-KEY.



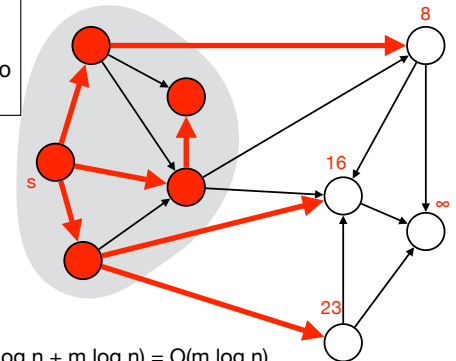
Dijkstra's Algorithm

```

DIJKSTRA(G, s)
  for all vertices v ∈ V
    v.d = ∞
    v.π = null
  INSERT(P, v)
  DECREASE-KEY(P, s, 0)
  while (P ≠ ∅)
    u = EXTRACT-MIN(P)
    for all v that u point to
      RELAX(u, v)
    
```

```

RELAX(u, v)
  if (v.d > u.d + w(u, v))
    v.d = u.d + w(u, v)
    DECREASE-KEY(P, v, v.d)
    v.π = u
    
```



- **Time.**
 - n EXTRACT-MIN
 - n INSERT
 - $< m$ DECREASE-KEY
- **Total time with min-heap.** $O(n \log n + n \log n + m \log n) = O(m \log n)$

Dijkstra's Algorithm

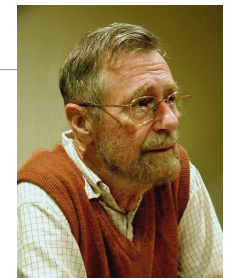
- **Priority queues and Dijkstra's algorithm.** Complexity of Dijkstra's algorithm depend on priority queue.
 - n INSERT
 - n EXTRACT-MIN
 - $< m$ DECREASE-KEY

Priority queue	INSERT	EXTRACT-MIN	DECREASE-KEY	Total
array	$O(1)$	$O(n)$	$O(1)$	$O(n^2)$
binary heap	$O(\log n)$	$O(\log n)$	$O(\log n)$	$O(m \log n)$
Fibonacci heap	$O(1)^\dagger$	$O(\log n)^\dagger$	$O(1)^\dagger$	$O(m + n \log n)$

$\dagger =$ amortized

- **Greedy.** Dijkstra's algorithm is a **greedy** algorithm.

Edsger W. Dijkstra



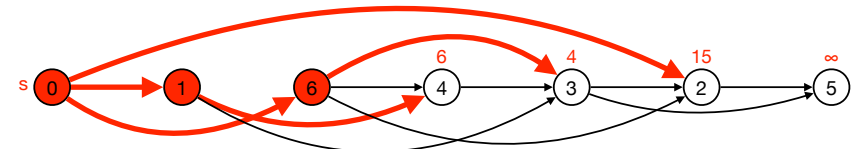
- Edsger Wybe Dijkstra (1930-2002)
- **Dijkstra algorithm.** "A note on two problems in connexion with graphs". Numerische Mathematik 1, 1959.
- **Contributions.** Foundations for programming, distributed computation, program verifications, etc.
- **Quotes.** "Object-oriented programming is an exceptionally bad idea which could only have originated in California."
- "The use of COBOL cripples the mind; its teaching should, therefore, be regarded as a criminal offence."
- "APL is a mistake, carried through to perfection. It is the language of the future for the programming techniques of the past: it creates a new generation of coding bums."

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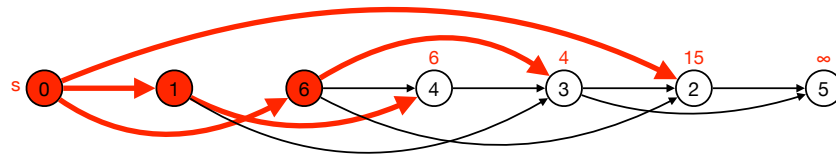
Shortest Paths on DAGs

- **Challenge.** Is it computationally easier to find shortest paths on DAGs?
- **DAG shortest path algorithm.**
 - Process vertices in topological order.
 - For each vertex v , relax all edges from v .
- Also works for **negative** edge weights.



Shortest Paths on DAGs

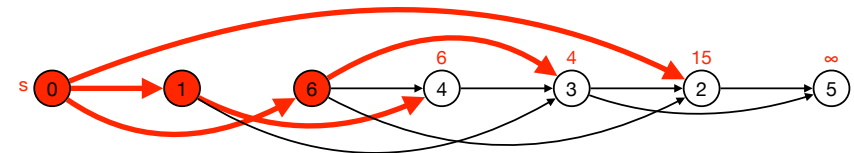
- **Lemma.** Algorithm computes shortest paths in DAGs.



- **Proof.**
 - Consider some step after growing tree T and assume distances in T are correct.
 - Consider next vertex u of s **not** in T .
 - Any path to u consists vertices in T + edge e to u .
 - Edge e is relaxed \implies distance to u is shortest.

Shortest Paths on DAGs

- **Implementation.**
 - Sort vertices in topological order.
 - Relax outgoing edges from each vertex.
- **Total time.** $O(m + n)$.



Shortest Paths Variants

- **Vertices**
 - Single source.
 - Single source, single target.
 - All-pairs.
- **Edge weights.**
 - Non-negative.
 - Arbitrary.
 - Euclidian distances.
- **Cycles.**
 - No cycles
 - No negative cycles.

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