

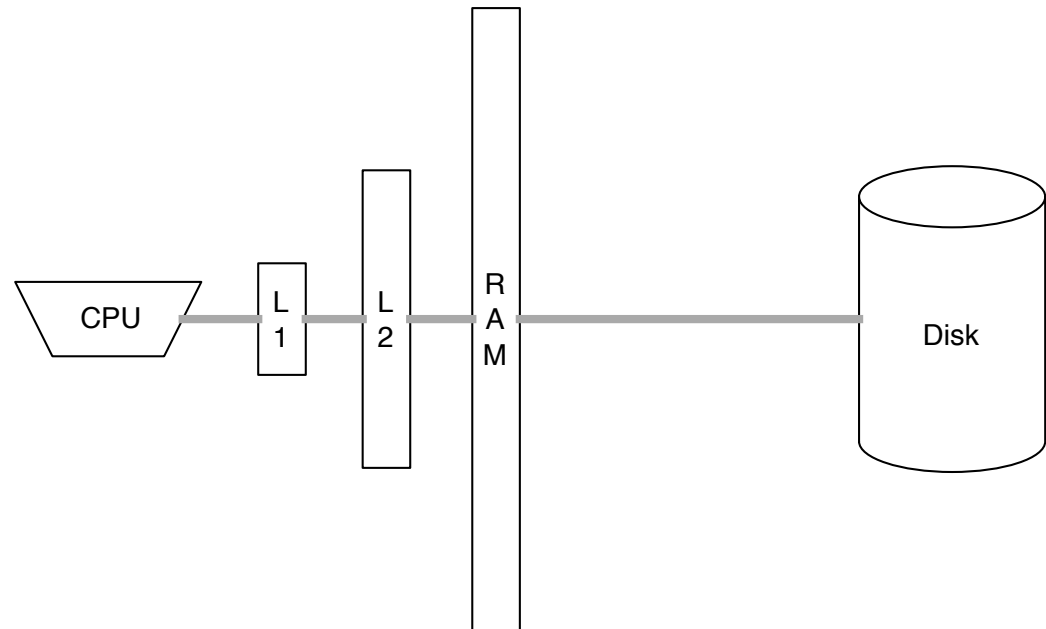
External Memory

- Computational models
- Shortest path in implicit grid graphs
 - RAM algorithm
 - I/O algorithms
 - Cache-oblivious algorithm

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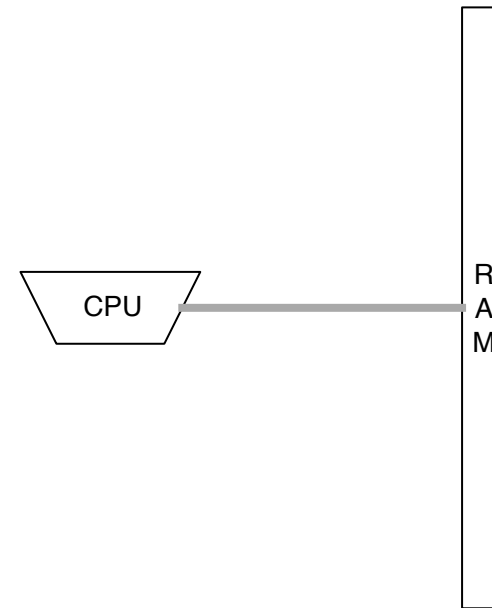
Computational Models



- iPad Air 2.

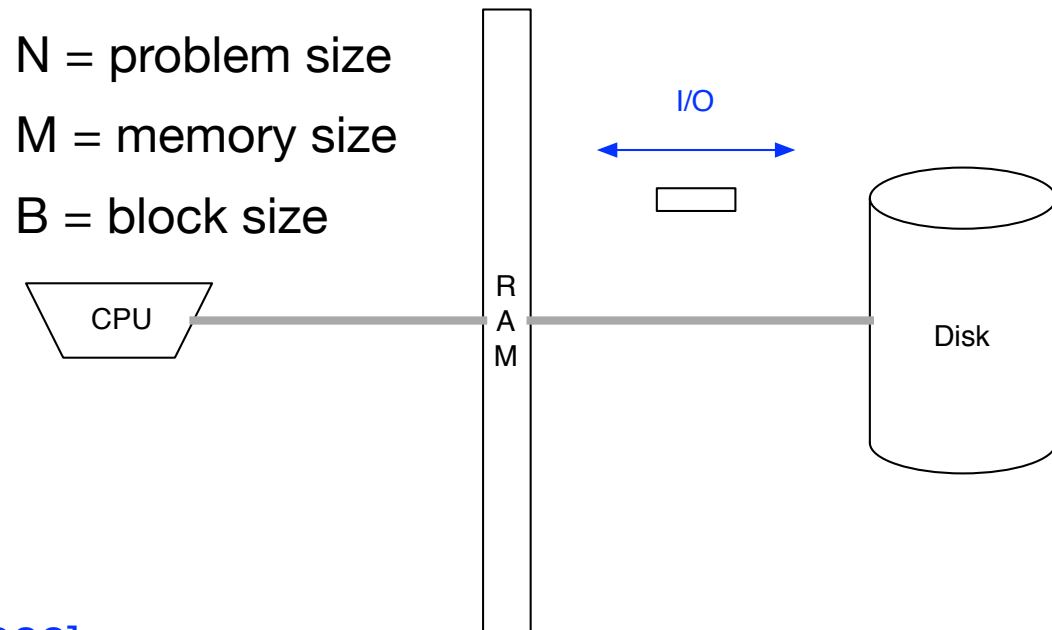
- A8X Chip (triple-core ARMv8-A)
- L1 cache: 64 KB instruction + 64 KB data per core
- L2 cache: 2 MB
- L3 cache: 4 MB
- Memory: 2 GB
- Disk: 16 GB SSD

Computational Models



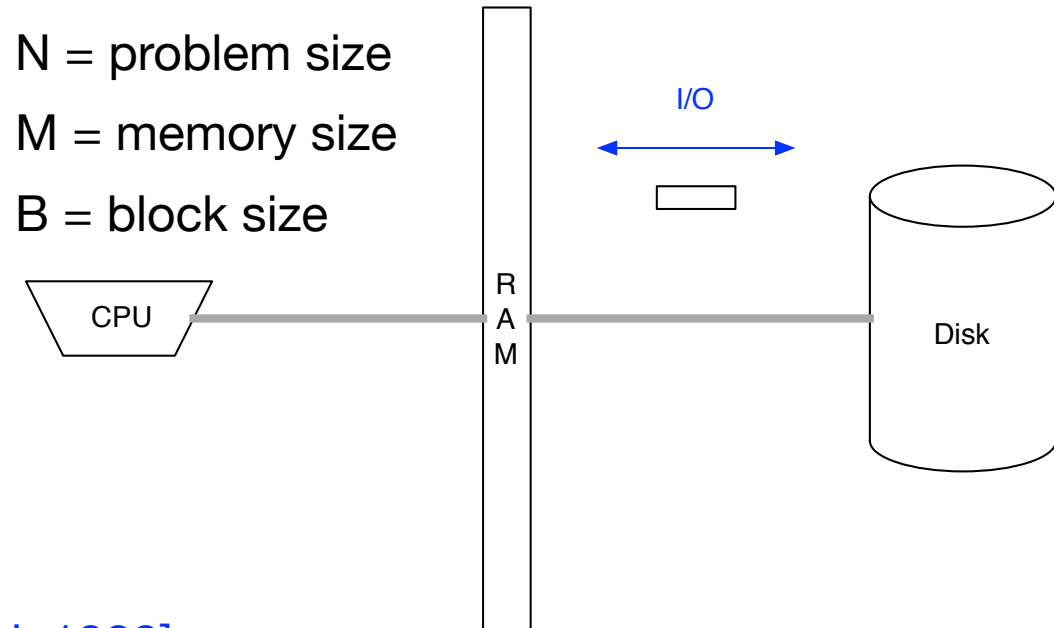
- Word RAM model.
 - Infinite memory of cells.
 - Read/write a cell.
 - Arithmetic and boolean operations (+, -, /, =, <, >, &, |, ...)
- Cost.
 - Time complexity = number of operations.

Computational Models



- I/O model [Aggarwal and Vitter 1988].
 - Limited memory + infinite **disk**
 - I/O operation = read/write consecutive **block** of B cells between memory and disk.
 - Arithmetic and boolean operations ($+$, $-$, $/$, $=$, $<$, $>$, $\&$, $|$, ...) on cells in memory.
- **Cost.**
 - I/Os = number of I/O operations.
 - Computation is free (!)

Computational Models

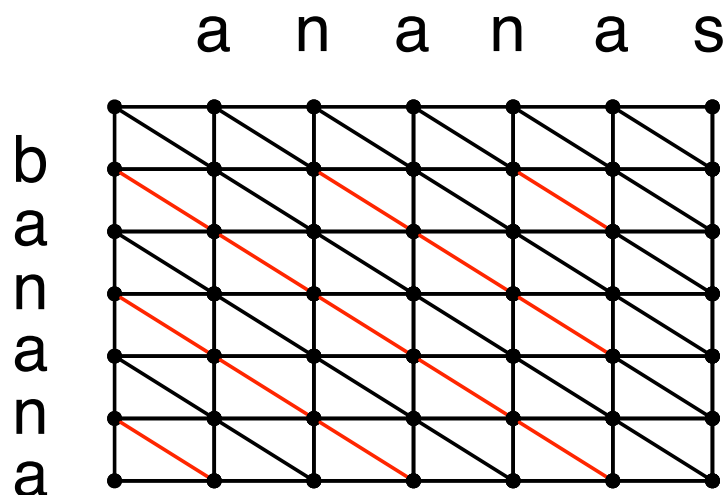


- Cache-oblivious model [Frigo et al. 1999].
 - Identical to I/O model except algorithms do not know M and B .
 - Program in RAM model and analyze in I/O model.
 - Assume optimal cache replacement strategy with full associativity.
- Properties.
 - Efficient on one level of cache \Rightarrow efficient on all levels cache.
 - Portable + self-tunable + simple.

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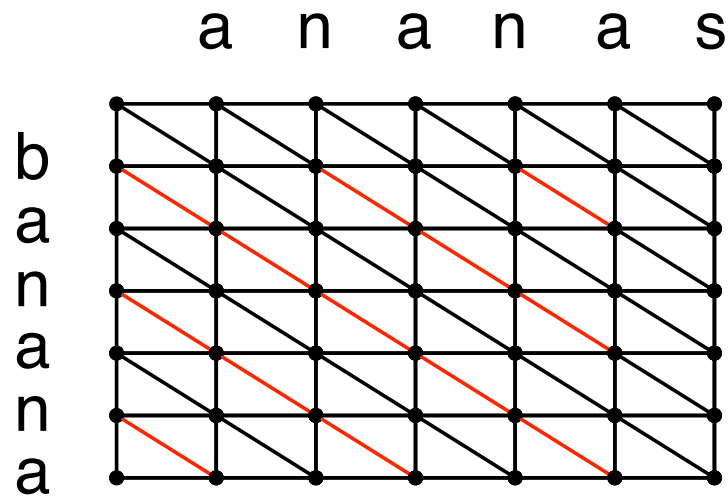
Shortest Paths in Implicit Grid Graphs



- **Implicit grid graphs.**

- Let S and T be strings of length n .
- The **implicit grid graph** for S and T is A 2D grid of $(n+1) \times (n+1)$ nodes.
 - For each node an edge to neighbors to E, S, SE.
 - E and S edges have weight 1.
 - SE edge $(i-1, j-1)$ to (i, j) has weight 0 if $S[i] = T[j]$ and 1 otherwise.

Shortest Paths in Implicit Grid Graphs



- Shortest paths in implicit grid graphs (SPIIG) problem.
 - Input. Strings S and T of length n.
 - Output. Length of shortest path from (0,0) to (n,n).

Shortest Paths in Implicit Grid Graphs

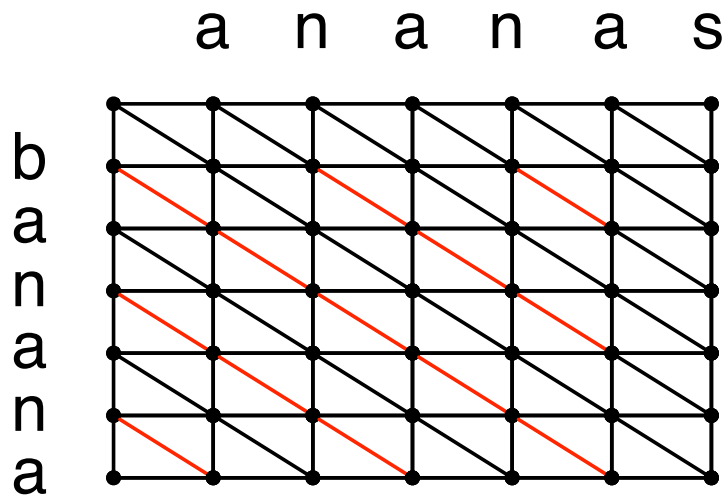
- Applications.

- Shortest paths in implicit grid graphs is the **edit distance** problem.
- With other edge weight functions we get longest common subsequence, sequence alignment, string similarity, approximate string matching, etc.

External Memory

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RAM Algorithm



a n a n a s

| | | | | | | | |
|---|---|---|---|---|---|---|---|
| b | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
| a | 1 | 1 | 2 | 3 | 4 | 5 | 6 |
| n | 2 | 1 | 2 | 2 | 3 | 4 | 5 |
| a | 3 | 2 | 1 | 2 | 2 | 3 | 4 |
| n | 4 | 3 | 2 | 1 | 2 | 2 | 3 |
| a | 5 | 4 | 3 | 2 | 1 | 2 | 3 |
| a | 6 | 5 | 4 | 3 | 2 | 1 | 2 |

- How can we solve SPIIGG on a RAM?
- **Dynamic programming algorithm.**
 - Construct $(n+1) \times (n+1)$ matrix.
 - Fill in each entry in $O(1)$ time in left-to-right top-to-bottom order.
- **Time.** $O(n^2)$
- **Space.** $O(n)$ (only store current + last row)
- Slightly faster solutions known [MP1980, Myers1999, CLZ2002, BFC2008]

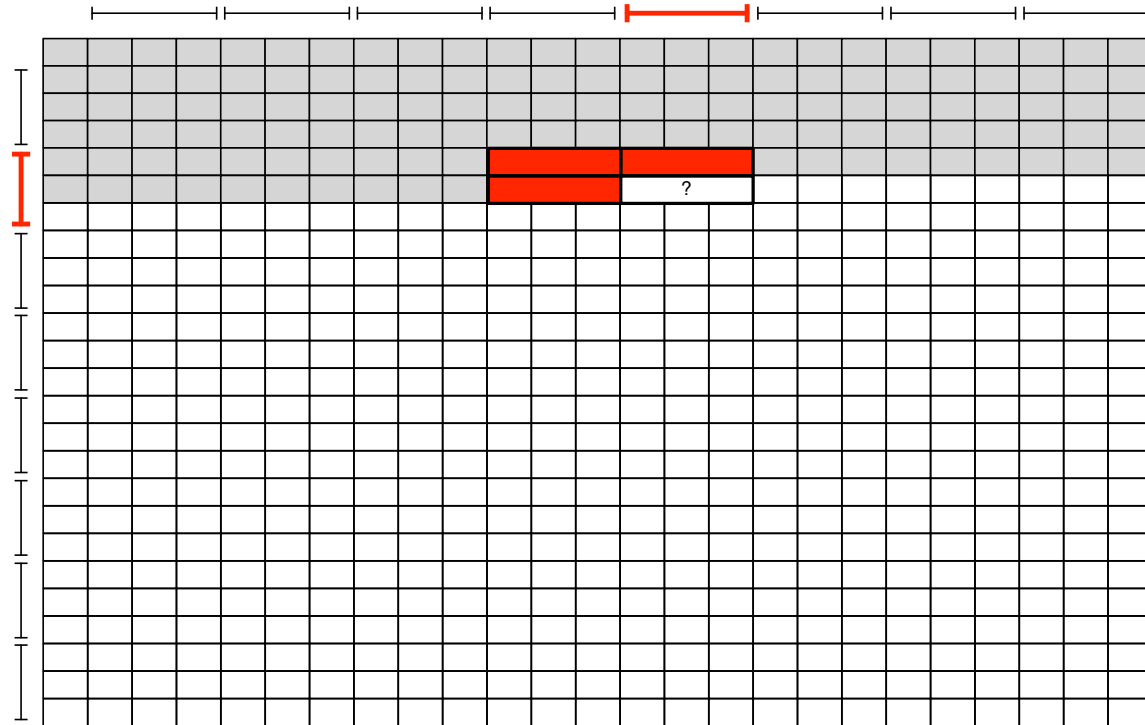
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External Memory Algorithms

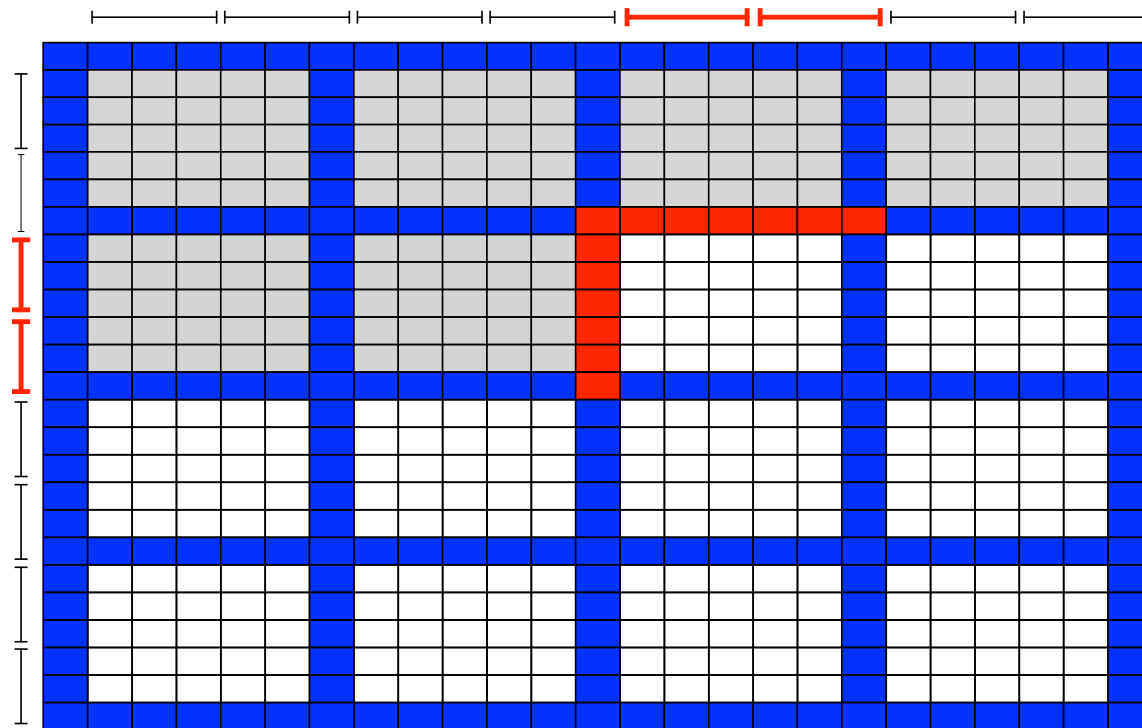
- **Goal.** Efficient external memory algorithms.
- **I/O model.**
 - **Solution 1.** Converted RAM algorithm
 - **Solution 2.** Table partitioning
- **Cache-oblivious model.**
 - **Solution 3.** Recursive table partitioning

Solution 1. Converted RAM Algorithm



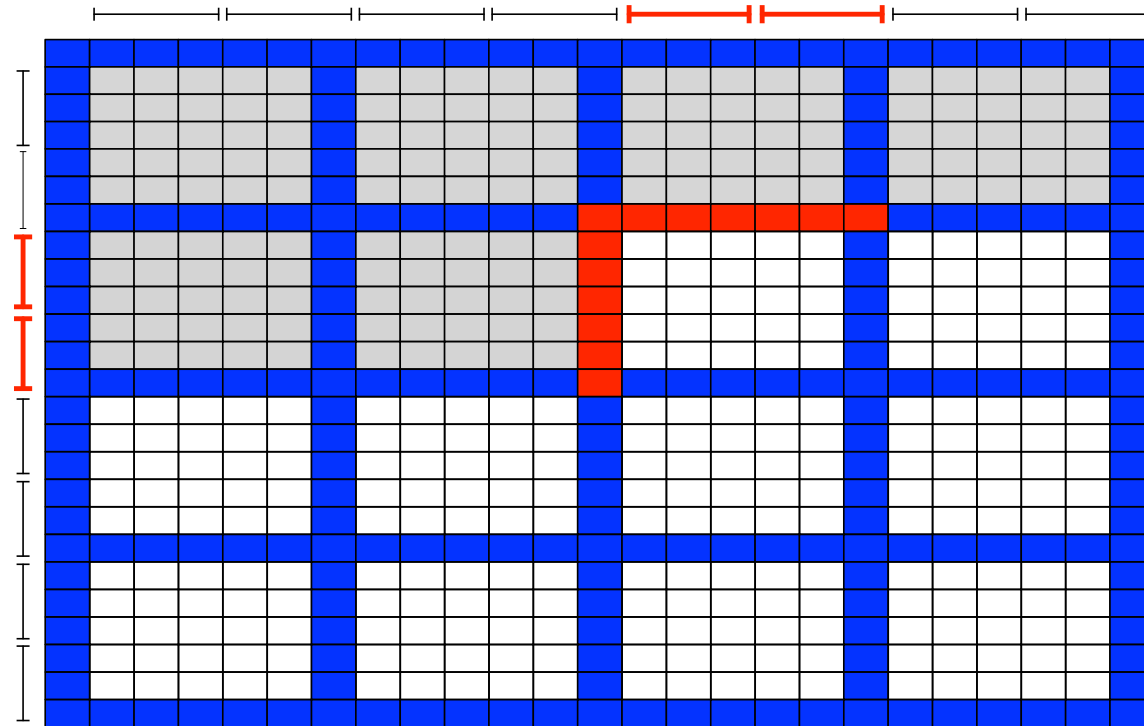
- Strings S and T stored consecutively in n/B blocks on disk.
- **Algorithm.**
 - Do as RAM algorithm. Read and write blocks as necessary.
- **I/Os.** $O(n^2/B)$.

Solution 2. Table Partitioning



- Divide into subtables with overlapping boundaries.
- **Algorithm.** Process subtables from left-to-right, top-to-bottom order. For each subtable:
 - Read corresponding substrings and input boundary into internal memory
 - Fill in subtable using RAM algorithm.
 - Write output boundary to disk.

Solution 2. Table Partitioning

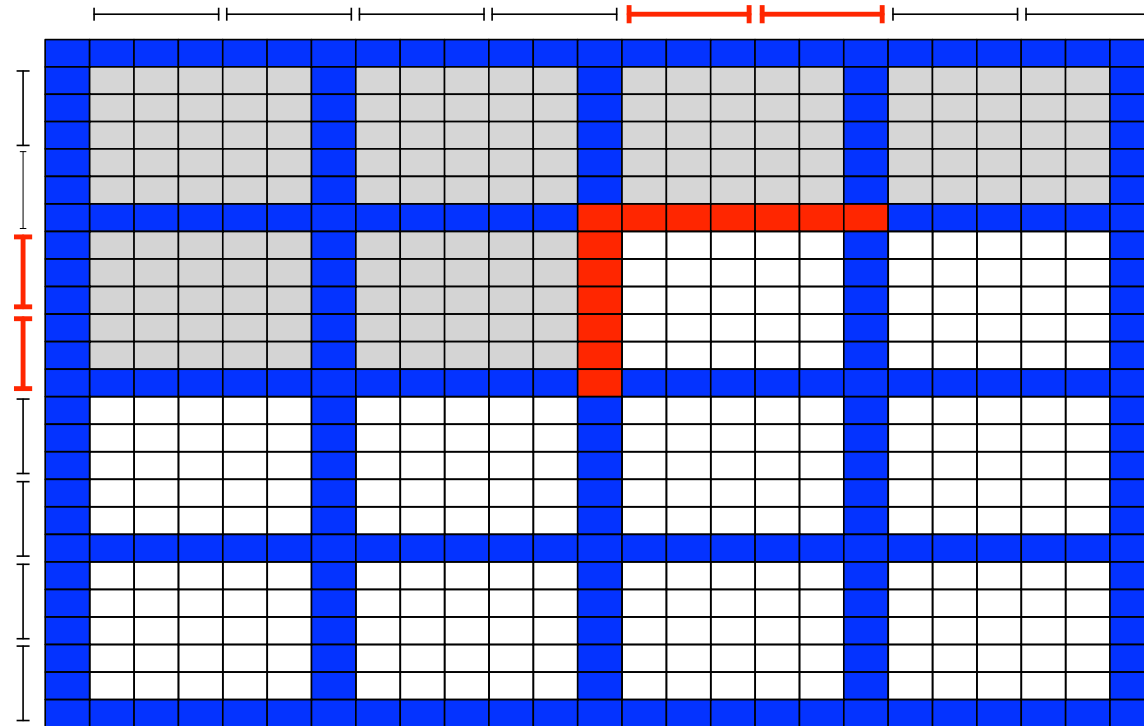


- How to choose subtable size?
- Make subtable $dM \times dM$ for $d < 1$ such that substrings + input boundary + output boundary + space for internal memory algorithm on subtable $< M$.
- I/Os.
 - Number of subtables = $O(n^2/M^2)$.
 - I/Os per subtable = $O(M/B)$.
 - $\Rightarrow O(n^2/M^2 \cdot M/B) = O(n^2/MB)$

Solution 2. Table Partitioning

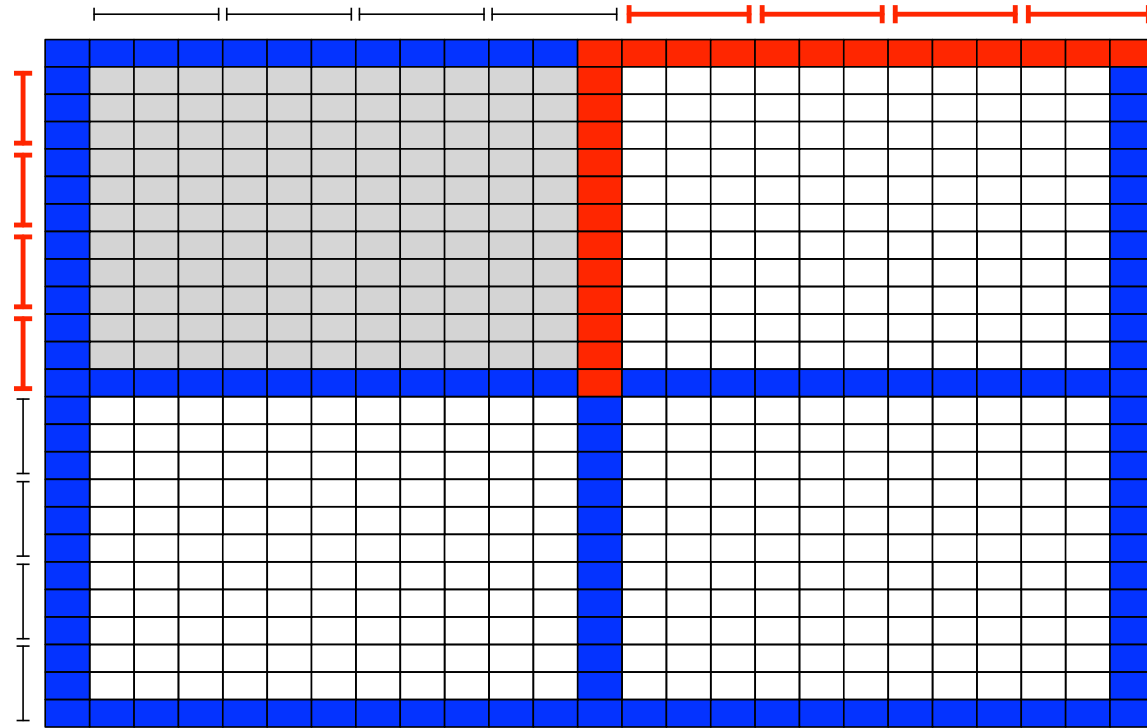
- **Theorem.** We can solve SPIIGG in the I/O model in
 - $O(n^2/MB + n/B)$ I/Os
 - $O(n^2)$ time
 - $O(n)$ space

Solution 2. Table Partitioning



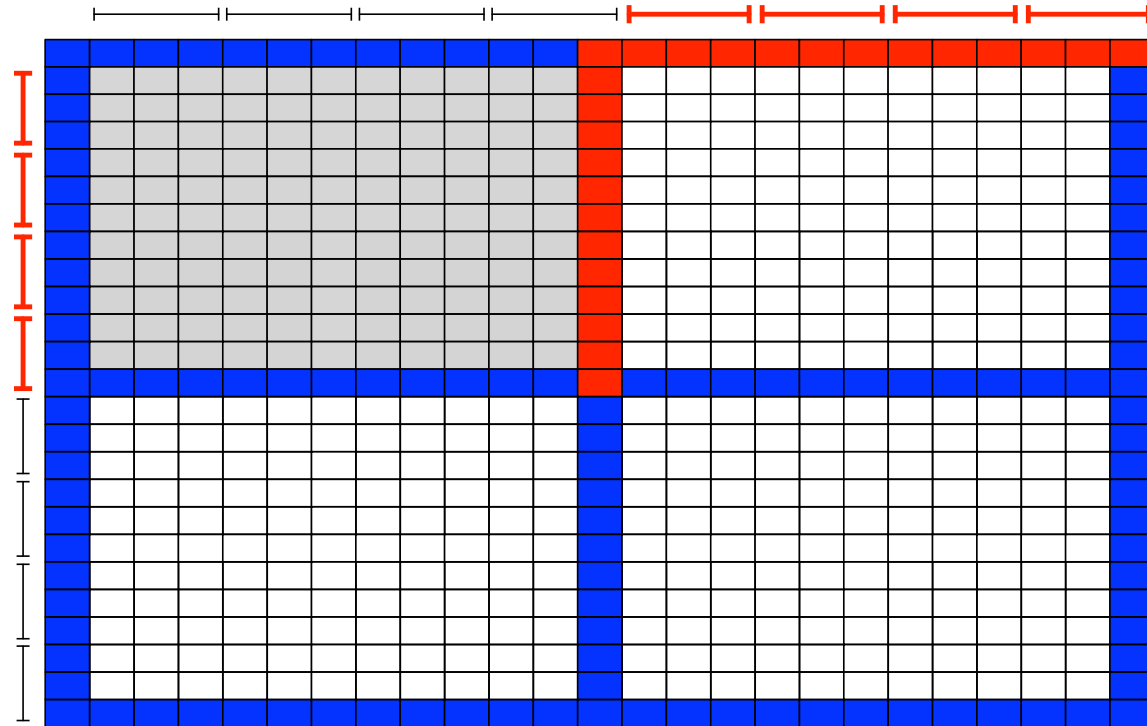
- How can we make solution 2 cache-oblivious?
- **Challenge.** We cannot use M and B .
- **Idea.** Use recursion to design algorithm that is good for **all** M and B .

Solution 3. Recursive Table Partitioning



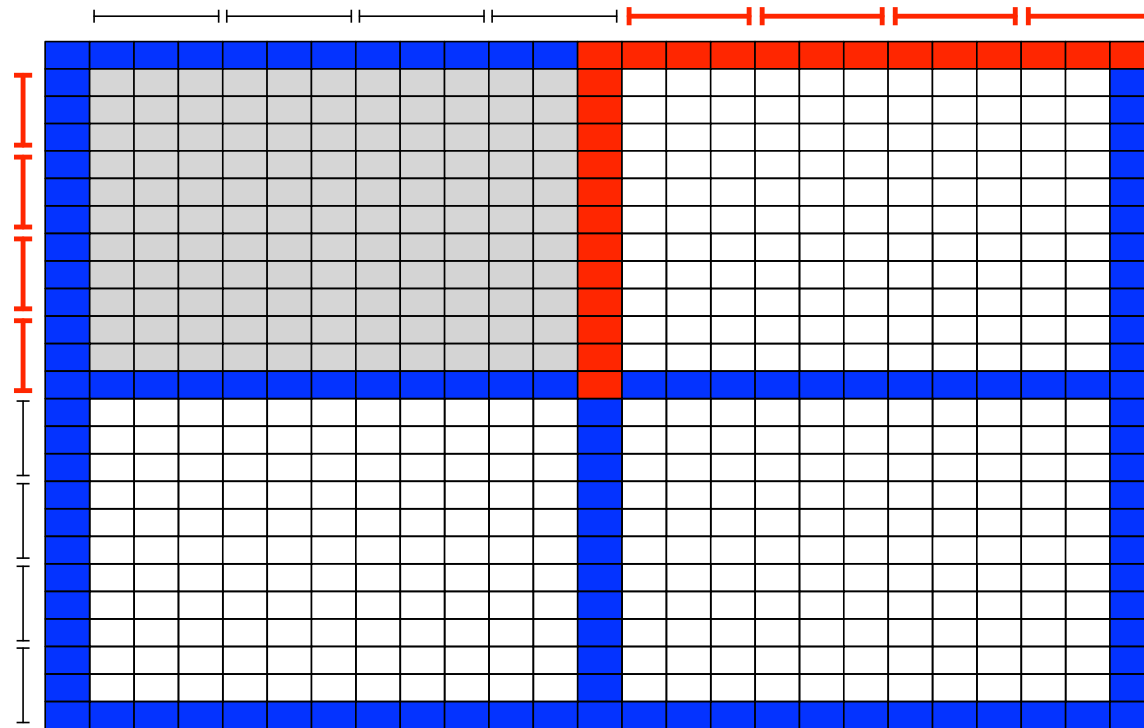
- [Algorithm.](#)
- Divide table into 4 quadrants with overlapping boundaries.
- Process quadrants from left-to-right, top-to-bottom order. For each quadrant:
 - Read corresponding substrings and input boundary.
 - Fill in quadrant recursively.
 - Write output boundary.

Solution 3. Recursive Table Partitioning



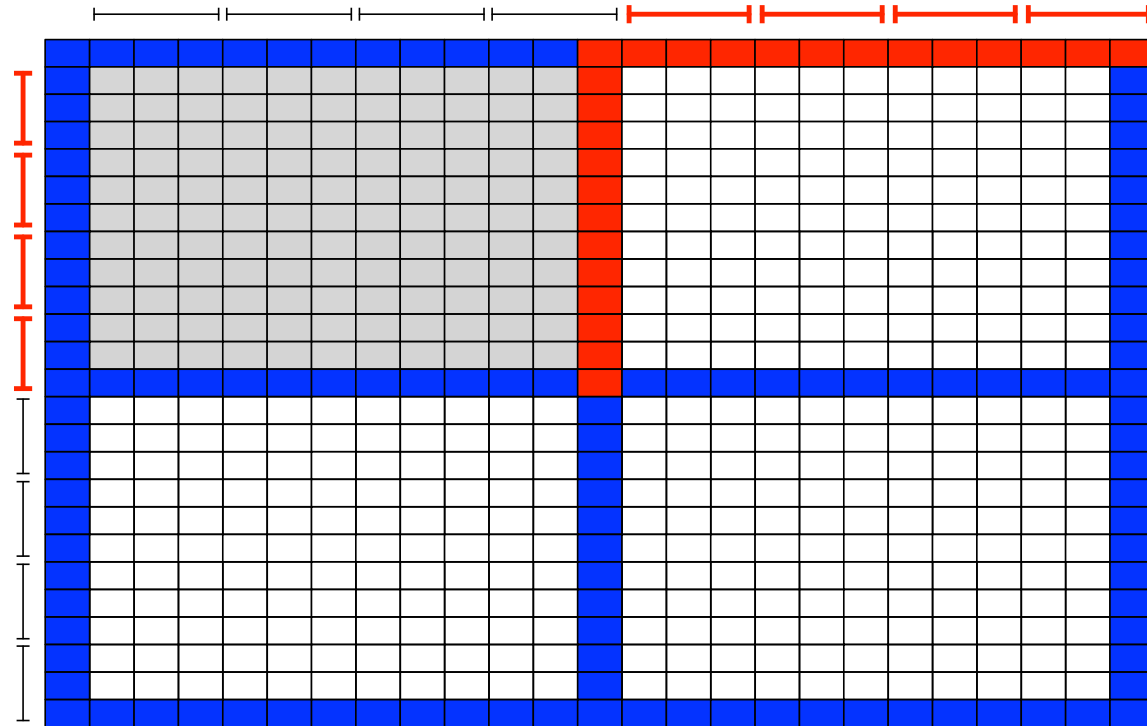
- I/Os.
- Define $IO(n)$ = number of I/Os to process a table of size $n \times n$
 - Case 1: $n \leq dM$ (substrings + boundaries + computation fit in internal mem)
 - $IO(n) = O(n/B)$
 - Case 2: $n > dM$?

Solution 3. Recursive Table Partitioning



- **Algorithm.**
- Divide table into 4 quadrants with overlapping boundaries. $O(1)$
- Process quadrants from left-to-right, top-to-bottom order. For each quadrant:
 - Read corresponding substrings and input boundary. $O(n/B)$
 - Fill in quadrant recursively. $4 \cdot IO(n/2)$
 - Write output boundary. $O(n/B)$

Solution 3. Recursive Table Partitioning



- Case 1 + 2:

$$IO(n) = \begin{cases} O(n/B) & \text{if } n \leq dM \\ 4 \cdot IO(n/2) + O(n/B) & \text{if } n > dM \end{cases}$$

- $\Rightarrow IO(n) = O(n^2/MB)$

Solution 3. Recursive Table Partitioning

- **Theorem.** We can solve SPIIGG in the cache-oblivious model in
 - $O(n^2/MB + n/B)$ I/Os
 - $O(n^2)$ time
 - $O(n)$ space

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