

Radix and Suffix Sorting

- Radix Sort
- Suffix Sort

Philip Bille

Radix and Suffix Sorting

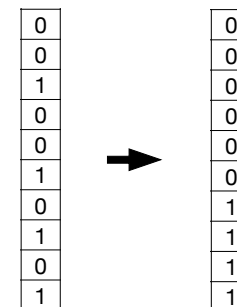
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Radix Sort

- [Sorting small universes](#). Given a sequence of n integers from a universe $U = \{0, 1, \dots, u-1\}$.
- How fast can we sort sequence if the size of the universe is not too big?

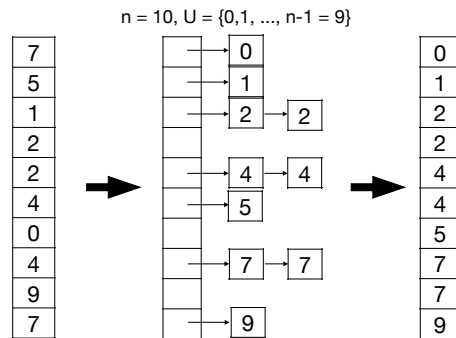
Radix Sort

$n = 10$, $U = \{0, 1\}$



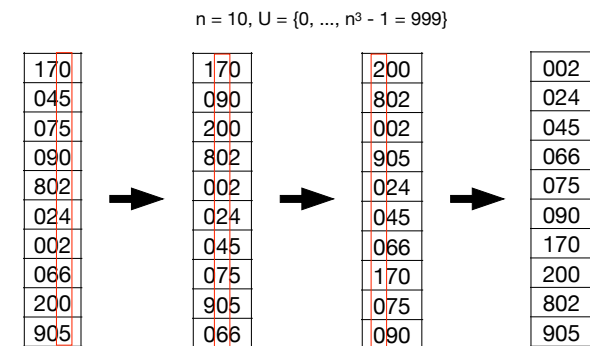
- [Algorithm](#). Count 0s and 1s.
- [Time](#). $O(n)$.

Radix Sort



- **Algorithm.** Insert into array of linked list + traverse array of linked list.
- **Time.** $O(n + u) = O(n)$
- Sorting can be **stable**.

Radix Sort



- **Radix Sort.** Sort on each digit from right to left using stable sort.
- **Time.** $O(n + n + n) = O(n)$

Radix Sort

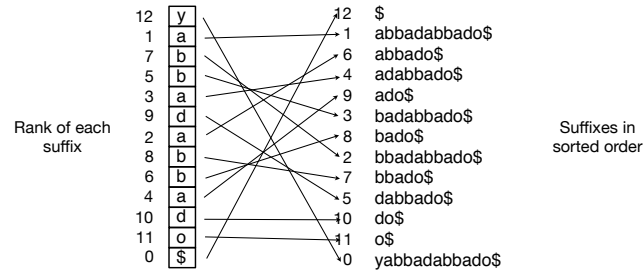
- **Radix Sort [Hollerith 1887].** Sort sequence of n integers from $U = \{0, \dots, n^k - 1\}$.
 - Write each element in sequence as a base n integer $x = (x_1, x_2, \dots, x_k)$
 - Sort sequence according to each digit from right to left. Sorting should be **stable**.
- **Time.** $O(nk)$

Radix and Suffix Sorting

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Suffix Sort

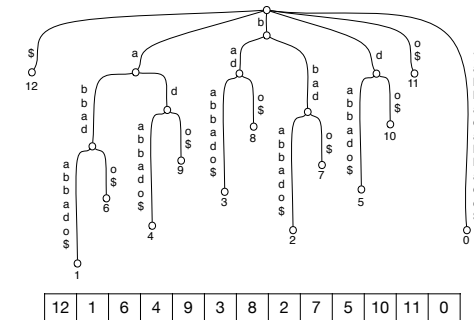
- **Suffix sorting.** Given string S of length n over alphabet Σ , compute the sorted **lexicographic order** of all suffixes of S .



- **Theorem [Kasai et al. 2001].** Given the sorted lexicographic order of suffixes of S , we can construct the suffix tree for S in linear time.

Suffix Sort

- **Suffix trees and sorting.** The lexicographic order of the suffixes is the same ordering as suffixes in the leaves of the suffix tree.
- **Suffix array.** The array of the sorted order of the suffixes.



Suffix Sort

- **Goal.** Compute the lexicographic order of all suffixes of S fast.
- For simplicity assume $|\Sigma| = O(n)$
- **Solution in 3 steps.**
 - Solution 1: Radix sorting
 - Solution 2: Prefix doubling
 - Solution 3: Difference cover sampling

Solution 1: Radix Sort

- **Radix Sort.**
 - Generate all suffixes (pad with \$).
 - Radix sort.

```

yabbadabbado$
abbadabbado$
bbadabbado$
badabbado$
adabbado$
dabbado$
abbado$
bbado$
bado$
ado$
do$
o$
$
    
```

- **Time.** $O(n^2)$

Solution 2: Prefix Doubling

- **Prefix doubling** [Manber and Myers 1990]. Sort substrings (padded with \$) of lengths 1, 2, 4, 8, ..., n. Each step uses radix sort on pair from previous step.

5	y	8	51	ya	10	84	yabb
1	a	1	12	ab	1	13	abba
2	b	4	22	bb	6	42	bbad
2	b	3	21	ba	4	35	bada
1	a	2	13	ad	2	21	adab
3	d	5	31	da	7	54	dabb
1	a	1	12	ab	1	13	abba
2	b	4	22	bb	6	42	bbad
2	b	3	21	ba	5	36	bado
1	a	2	13	ad	3	27	ado\$
3	d	6	34	do	8	60	do\$\$
4	o	7	40	o\$	9	70	o\$\$\$
0	\$	0	00	\$\$	0	00	\$\$\$\$

- **Time.** $O(n \log n)$

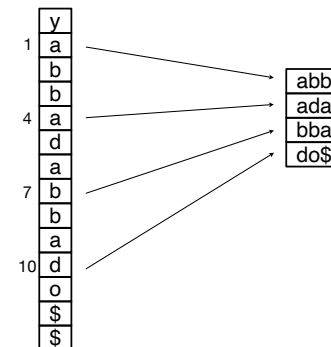
Solution 3: Difference Cover Sampling

- **DC3 Algorithm** [Karkkainen et al. 2003]. Sort suffixes in three steps:
 - **Step 1.** Sort sample suffixes.
 - Sample all suffixes starting at positions $i = 1 \bmod 3$ and $i = 2 \bmod 3$.
 - Recursively sort sample suffixes.
 - **Step 2.** Sort non-sample suffixes.
 - Sort the remaining suffixes (starting at positions $i = 0 \bmod 3$).
 - **Step 3.** Merge.
 - Merge sample and non-sample suffixes.

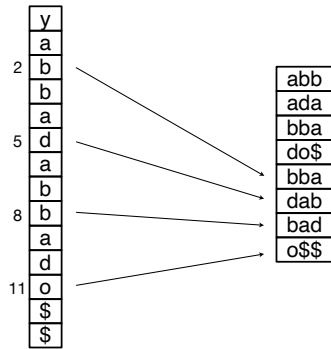
Step 1: Sort Sample Suffixes

y
a
b
b
a
d
a
b
b
a
d
o
\$
\$

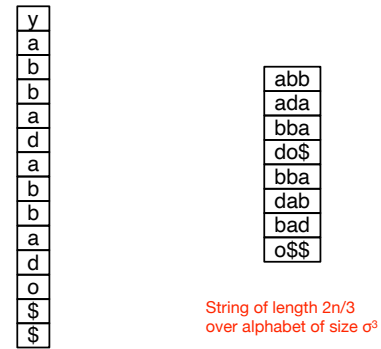
Step 1: Sort Sample Suffixes



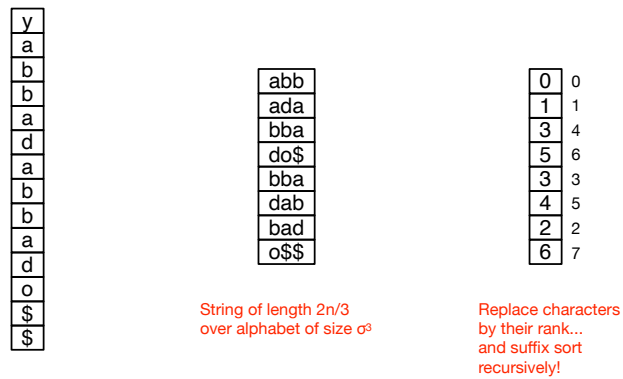
Step 1: Sort Sample Suffixes



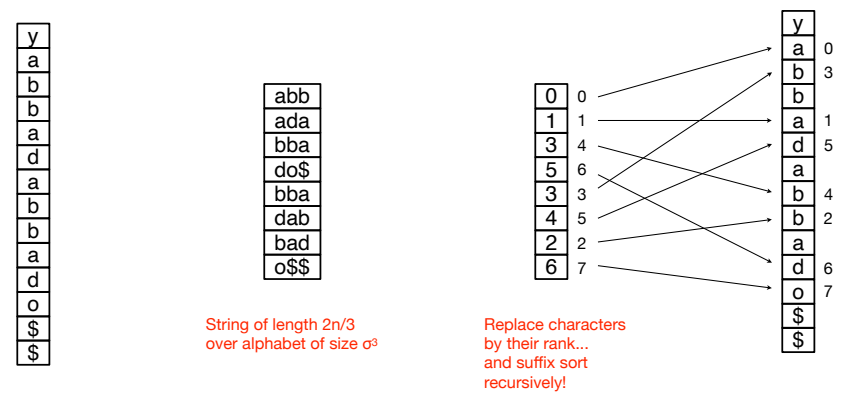
Step 1: Sort Sample Suffixes



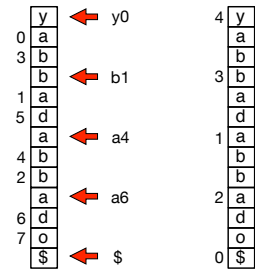
Step 1: Sort Sample Suffixes



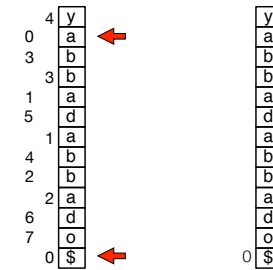
Step 1: Sort Sample Suffixes



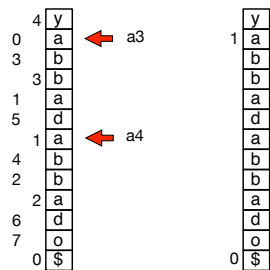
Step 2: Sort Non-Sample Suffixes



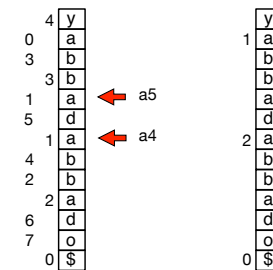
Step 3: Merge



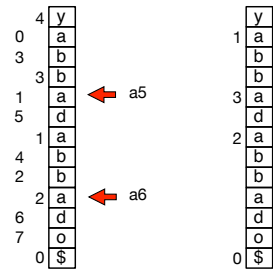
Step 3: Merge



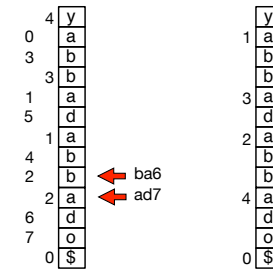
Step 3: Merge



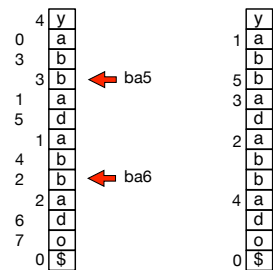
Step 3: Merge



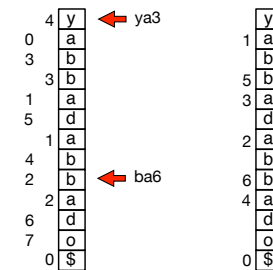
Step 3: Merge



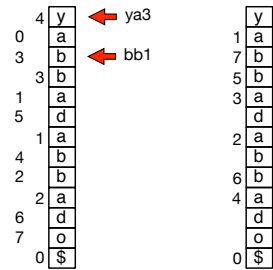
Step 3: Merge



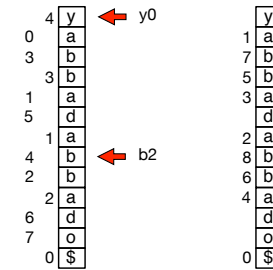
Step 3: Merge



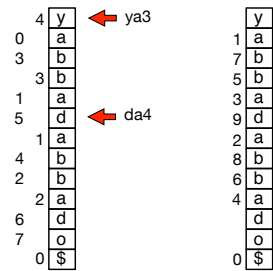
Step 3: Merge



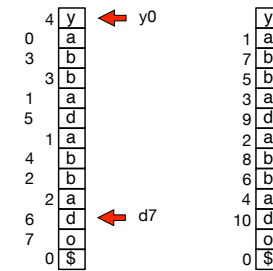
Step 3: Merge



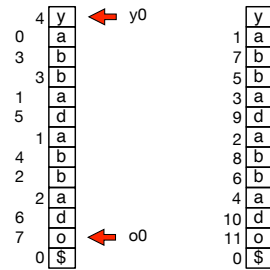
Step 3: Merge



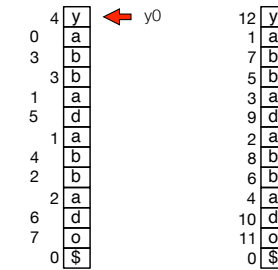
Step 3: Merge



Step 3: Merge

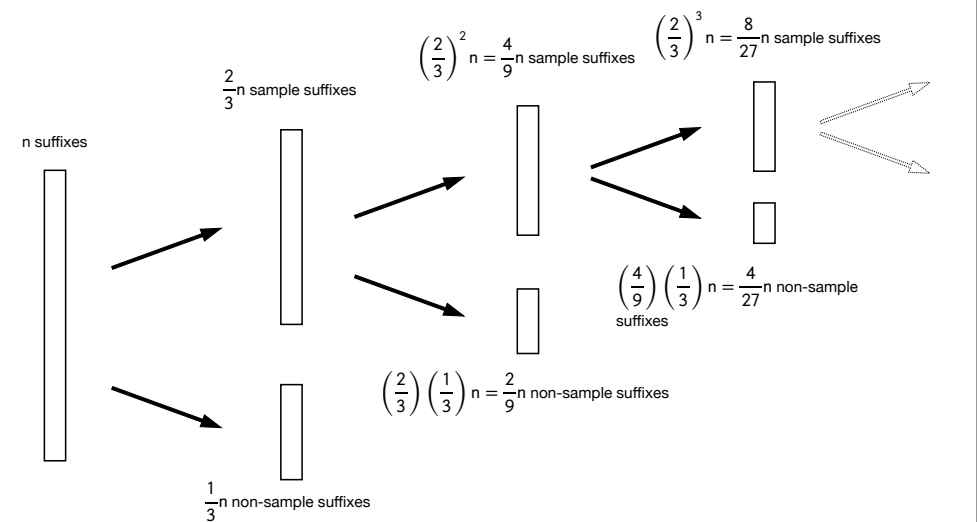


Step 3: Merge



Solution 3: Difference Cover Sampling

- **DC3 Algorithm.** Sort suffixes in three steps:
 - **Step 1.** Sort sample suffixes.
 - Sample all suffixes starting at positions $i = 1 \bmod 3$ and $i = 2 \bmod 3$. $O(n)$
 - Recursively sort sample suffixes. $T(2n/3)$
 - **Step 2.** Sort non-sample suffixes.
 - Sort the remaining suffixes (starting at positions $i = 0 \bmod 3$). $O(n)$
 - **Step 3.** Merge.
 - Merge sample and non-sample suffixes. $O(n)$
- **Time.** $T(n) = T(2n/3) + O(n) = O(n)$



Solution 3: Difference Cover Sampling

- **Theorem.** We can suffix sort a string of length n over alphabet Σ of size n in time $O(n)$.
- **Theorem.** We can suffix sort a string of length n over alphabet Σ $O(\text{sort}(n, |\Sigma|))$ time.

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