## Weekplan: External Memory I

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## **References and Reading**

- [1] The Input/Output Complexity of Sorting and Related Problems, A. Aggarwal and J. Vitter, CACM 1988. Set P = 1 when reading this.
- [2] Organization and Maintenance of Large Ordered Indexes, R. Bayer, E. McCreight, Acta Inform., 1972.
- [3] Introduction to Algorithms, 3rd edition, Chap. 18, T.H. Cormen, C.E. Leiserson, R.L. Rivest, C. Stein, 2009.

We recommend reading [1] and [3] in detail. [2] is the original paper introducing *B*-trees.

## Exercises

**1** [*w*] **Prefix Sum** Given an array *A* of *N* elements, the *prefix-sum* of *A* is the array *P* such that  $P[i] = \sum_{j \le i} A[i]$ . Show how to compute the prefix sum of *A* efficiently in external memory

**2** [*w*] **Memory Hierarchy** Determine the configuration of the memory hierarchy on your own computer. Also, what is the cache-inclusion policy?

- 3 Stacks and Queues Consider stacks and queue in external memory. Solve the following exercises.
- **3.1** Show how to efficiently implement a stack in external memory. What is the worst-case and amortized I/Os per operation?
- **3.2** Do the same for a queue.

**4** Linked Lists Consider a data structure that maintains a sequence of element  $L = e_1, \ldots, e_N$  under the following operations:

- insert(e, e'): Insert element e' immediately after element e in the sequential order in L (extending the length of the sequence by 1).
- delete(*e*): Delete the element *e* in *L*.
- traverse(): Report the elements in *L* in sequence.

We assume that the argument e is a pointer to the element. Show how to efficiently implement the operations in external memory. *Hint:* What are the optimal I/O bounds for these operations you can hope to achieve? Try to achieve that.

**5 Dictionaries and Hashing** Recall the standard dictionaries based on hashing from your previous algorithms studies. What are the I/O complexities for these solutions? Consider the *exact* number of I/Os.

**6 Range Reporting** Suppose we want to extend *B*-trees to support the following range reporting operation:

• report(*i*, *j*): Report all elements with keys *k*, such that  $i \le k \le j$ .

Show how to efficiently implement report on *B*-trees. Your solution should have a good dependency on the size of the output.

**7 Insertions in** *B***-tree** Consider the following *B*-tree of order 4. The capital letters represent subtrees. Show the the tree after inserting 59.



**8** *B*-tree Construction Show how to efficiently construct a *B*-tree from an array of *N* elements.

**9 Optimality of** *B***-trees** Suppose that we want to search among *N* keys. Furthermore, suppose that the only way of accessing disk blocks is by following pointers. Show that a search takes at least  $\Omega(\log_B N/M)$  I/Os in the worst case. *Hint:* Consider the size  $C_t$  of the set of blocks that can be accessed in at most *t* I/Os. Assume that our memory initially is full of pointers.

**10 Dynamic Programming** Let *S* and *T* be strings of length *N* and consider the classic  $O(N^2)$  time solution for computing the longest common subsequence of *S* and *T*. Show how to implement the algorithm efficiently in external memory.