Weekplan: Massively Parallel Computation

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References and Reading

- [1] Massively Parallel Algorithms, M. Ghaffari, 2019.
- [2] Parallel Algorithms, Chapter 6, M. Ghaffari, 2019.
- [3] Scribe notes from "Algorithms for Massive Data Science", Ben Mosely, 2019.

We recommend reading [1] for an overview of the area. We recommend reading [2] sections 6.7-6.9 and the scribe notes [3] in detail.

Exercises

1 Summing with other Parameters Consider the sum problem from the lectures and the parameters *P* and *S*. Solve the following exercises.

- **1.1** Give an efficient algorithm when $S = \Theta(N^{3/4})$ and $P = \Theta(N^{1/4})$.
- **1.2** Give an efficient algorithm when $S = \Theta(N^{1/4})$ and $P = \Theta(N^{3/4})$.
- **1.3** Give an efficient algorithm when $S = \Theta(N^{\varepsilon})$ and $P = \Theta(N^{1-\varepsilon})$.
- 2 Sorting Properties and Output Solve the following exercises.
- **2.1** [*w*] The single round algorithm for sorting in the MPC model is often called "quicksort" or "generalized quicksort". Why? What is the connection to quicksort?
- **2.2** The output of the MPC sorting algorithm is not in the required format. Show how to convert it to the required format in O(1) additional rounds.
- 3 Prefix Sum, Distinct Elements, and Word Count Solve the following exercises. Assume $S = P = \widetilde{\Theta}(\sqrt{N})$
 - **3.1** Let *A* be an array of *N* integers distributed among processors. Each entry in *A* is stored as (i, A[i]). Show how to compute the prefix sum of *A* (the array $P[i] = \sum_{j \le i} A[j]$) efficiently in the MPC model. The array *P* should be represented similar to *A*.
 - 3.2 Let *L* be a list of *N* integers, show how to compute the number of distinct elements in *L* in the MPC model.
 - **3.3** Let *W* be a list of *N* strings each of constant length. The *word count* of *W* is the list of pairs of distinct words and their frequency in *W*. Computing the word count is a classic "hello world"-exercise for the MapReduce framework. Show how to implement it efficiently in the MPC model.
- 4 Sorting Analysis Solve the following exercises.
- **4.1** Carefully verify that each step of the sorting algorithm satisfies the bounds on *S* and *P*.
- **4.2** Show that $|X| \le 4P \ln N$ whp. *Hint:* First compute the expected size of X and then apply a Chernoff bound.

5 Minimum Spanning Tree Correctness Show that the MPC minimum spanning tree algorithm correctly outputs a minimum spanning tree. *Hint:* show that the discarded edges cannot be part of an MST using standard properties of MSTs.

6 Dynamic Programming Let *S* and *T* be strings of length *N* and consider the classic $O(N^2)$ time solution for computing the longest common subsequence of *S* and *T*. Show how to implement the algorithm efficiently on the MPC model. Assume $S = P = \tilde{\Theta}(\sqrt{N})$.

7 [*] **String Matching** Let *P* and *T* be strings of lengths *M* and *N*, respectively. The *string matching problem* is to determine if *P* occurs as a substring of *T*. Show how to solve the string matching problem efficiently on the MPC model. Assume $S = P = \tilde{\Theta}(\sqrt{N})$. Further, assume that *T* and *P* are partitioned into equal sized substrings of size *O*(*S*) and stored distributed among the processors. Solve the following exercises.

- 7.1 Solve the string matching problem with the assumption that M < S.
- **7.2** Solve the string matching problem for any $M \leq N$. *Hint*: use *Karp-Rabin fingerprints* to efficiently hash substrings.