

# Bit tricks and partial sums

Nicola Prezza

## References and Reading

- [1] Chapter 3 of: Navarro, Gonzalo. Compact data structures: A practical approach. Cambridge University Press, 2016.

## Exercises

**1 Bit tricks** The operation `popcount` takes as input a word  $X$  of  $w$  bits and outputs the number of 1s (i.e. bits set) in  $X$ . Assuming that  $w = 32$ , write the pseudocode of a function that computes `popcount(X)`. Tabulation (i.e. pre-computation of the answers in a table) is not allowed. The function should use only standard bitwise (`and`, `or`, `shift`) and arithmetic (addition, subtraction) operations. Try to use as few operations as possible (hint: for generic  $w$ , it is possible to use just  $O(\log w)$  instructions).

**2 PforDelta partial sum** The `PforDelta` code (see Chapter 2) encodes a sequence of  $n$  integers  $x_1, \dots, x_n$  as follows. We divide the numbers in consecutive blocks of size  $k$  (assume that  $k$  divides  $n$ ). Inside each block, we encode all integers using—for each of them—the bit-size of the largest integer in the block. The encoded  $x_1, \dots, x_n$  are then concatenated in a single packed bitvector  $B$ . Finally, we store the beginning of each block in  $B$  using  $n/k$  additional words.

Without loss of generality, assume that  $x_i \geq 1$ , for all  $1 \leq i \leq n$ , and let  $N = \sum_{i=1}^n x_i$ . Fix moreover the block size as  $k = w$ .

**2.1** Compute the bit-size of the `PforDelta` representation.

**2.2** What is the space overhead on top of the worst-case entropy of all integer sequences of length  $n$  that add up to  $N$ ?

**2.3** Show how to compute efficiently `access`<sup>1</sup>, `sum`, and `search` operations on top of the above representation.

---

<sup>1</sup>`access` simply outputs  $x_i$  given  $i$  as input